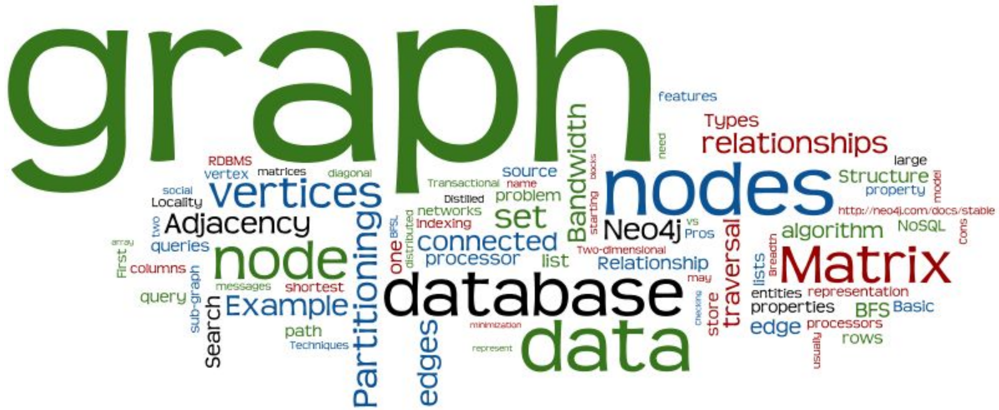


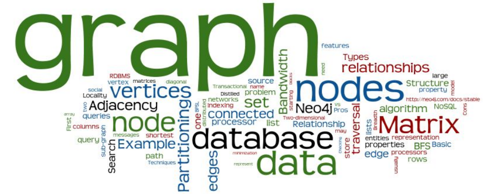
Graph Databases

Lecture 8 of *NoSQL Databases* (PA195)

David Novak, FI, Masaryk University, Brno

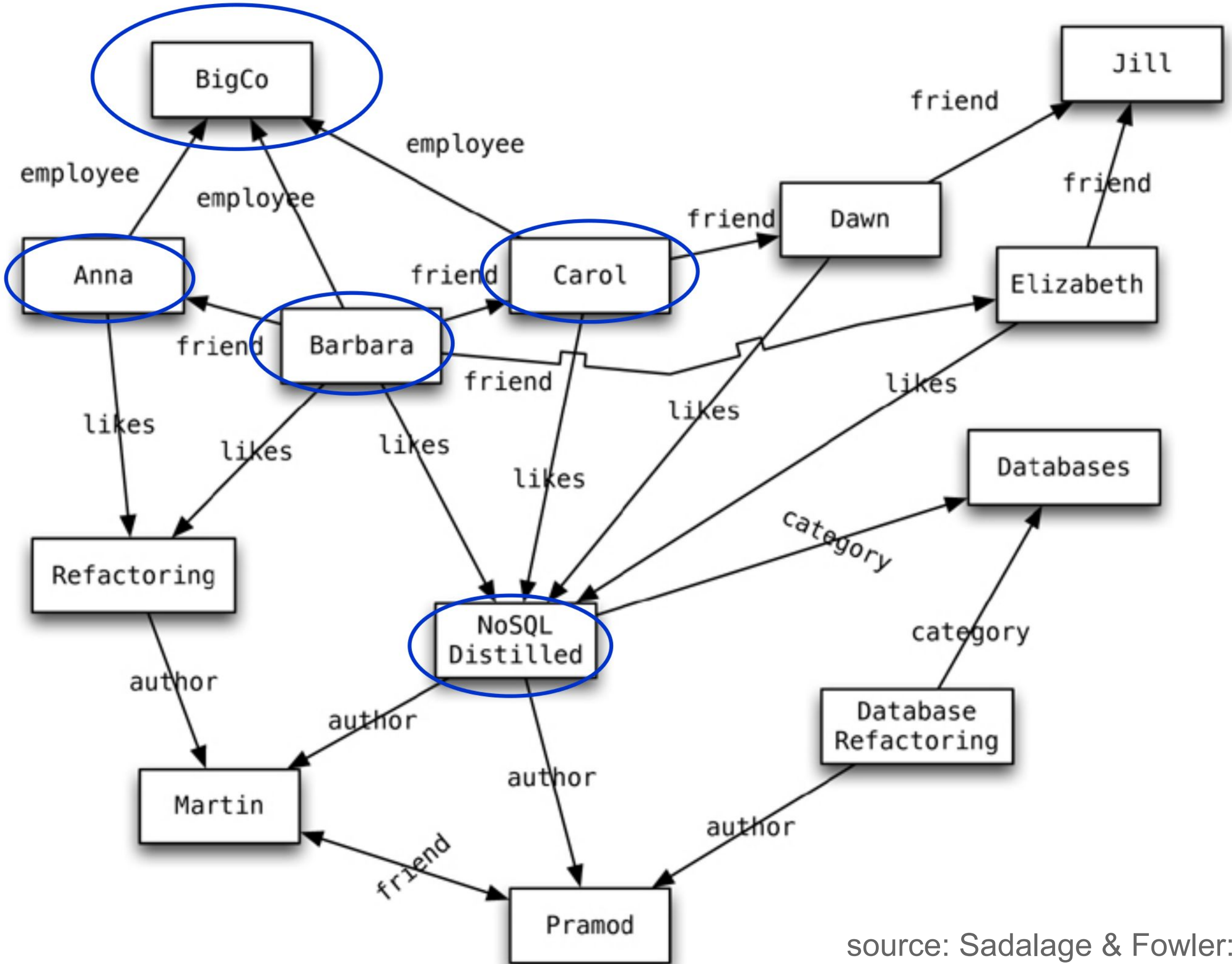
http://disa.fi.muni.cz/david-novak/teaching/nosql-databases-2016/

Agenda



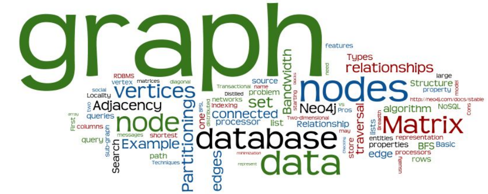
- Graph Databases: Mission, Data, Example
- A Bit of Graph Theory
  - Graph Representations
  - Algorithms: Improving Data Locality (efficient storage)
  - Graph Partitioning and Traversal Algorithms
- Graph Databases
  - Transactional databases
  - Non-transactional databases
- Neo4j
  - Basics, Native Java API, Cypher, Behind the Scene

Graph Databases: Example



source: Sadalage & Fowler: NoSQL Distilled, 2012

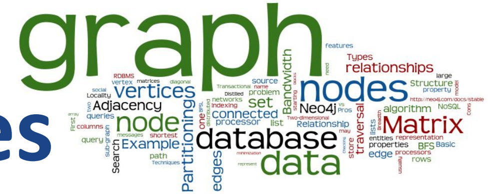
Graph Databases: Mission



- To store entities and relationships between them
  - Nodes are instances of objects
  - Nodes have properties, e.g., name
  - Edges connect nodes and have directional significance
  - Edges have types e.g., likes, friend, …

- Nodes are organized by relationships
  - Allow to find interesting patterns
  - **example:** Get all nodes that are “employee” of “Big Company” and that “likes” “NoSQL Distilled”

Graph Databases: Representatives



Neo4j

OrientDB®

APACHE GIRAPH

TITAN

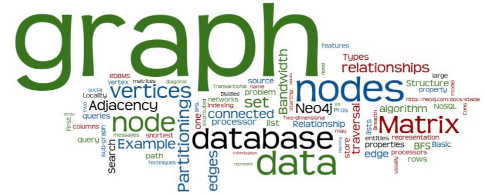
InfiniteGraph

Ranked list: http://db-engines.com/en/ranking/graph+dbms

A Bit of a Theory

Basics and graph representations

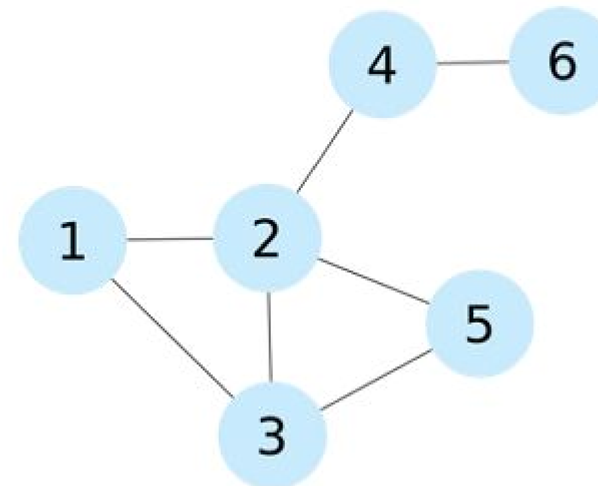
Basic Terminology



- Data: a set of entities and their relationships
  - => we need to efficiently represent graphs
- Basic operations:
  - finding the neighbours of a node,
  - checking if two nodes are connected by an edge,
  - updating the graph structure, …
  - => we need efficient graph operations
- Graph $G = (V, E)$ is usually modelled as
  - set of nodes (vertices) V, $|V| = n$
  - set of edges E, $|E| = m$
- Which data structure to use?

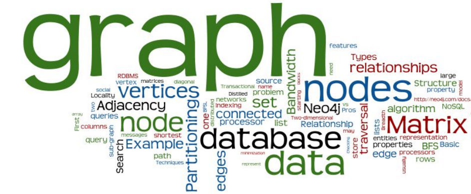
Data Structure: Adjacency Matrix

- Two-dimensional array A of $n \times n$ Boolean values
  - Indexes of the array = node identifiers of the graph
  - Boolean value A_{ij} indicates whether nodes i, j are connected

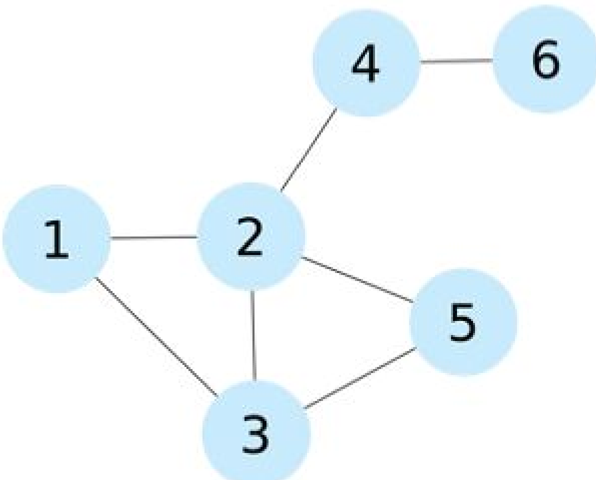


$$\begin{pmatrix} 0 & 1 & 1 & 0 & 0 & 0 \\ 1 & 0 & 1 & 1 & 1 & 0 \\ 1 & 1 & 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 & 0 & 1 \\ 0 & 1 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 \end{pmatrix}$$

- Variants:
  - (Un)directed graphs
  - Weighted graphs…



Adjacency Matrix: Properties

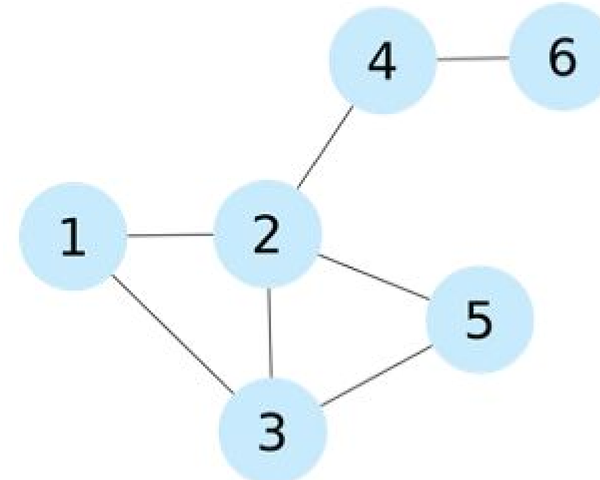


$$\begin{pmatrix} 0 & 1 & 1 & 0 & 0 & 0 \\ 1 & 0 & 1 & 1 & 1 & 0 \\ 1 & 1 & 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 & 0 & 1 \\ 0 & 1 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 \end{pmatrix}$$

- Pros:
  - Adding/removing edges
  - Checking if 2 nodes are connected
- Cons:
  - Quadratic space: $O(n^2)$
  - We usually have sparse graphs
  - Adding nodes is expensive
  - Retrieval of all the neighbouring nodes takes linear time: $O(n)$

Data Structure: Adjacency List

- A set of lists, each enumerating neighbours of one node
  - Vector of *n* pointers to adjacency lists
- Undirected graph:
  - An edge connects nodes *i* and *j*
  - => the adjacency list of *i* contains node *j* and vice versa
- Often compressed
  - Exploiting regularities in graphs



N1 → {N2, N3}

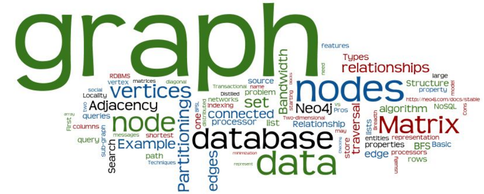
N2 → {N1, N3, N5}

N3 → {N1, N2, N5}

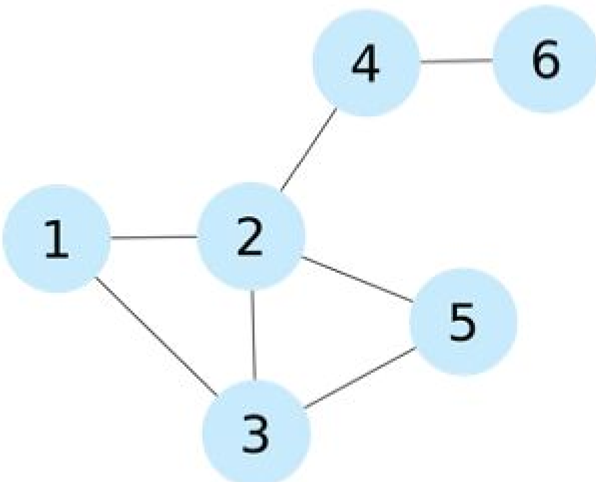
N4 → {N2, N6}

N5 → {N2, N3}

N6 → {N4}



Adjacency List: Properties



N1 → {N2, N3}

N2 → {N1, N3, N5}

N3 → {N1, N2, N5}

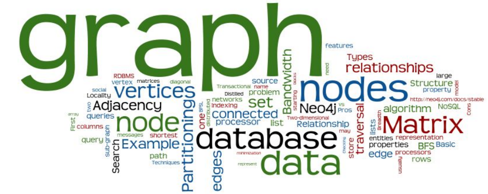
N4 → {N2, N6}

N5 → {N2, N3}

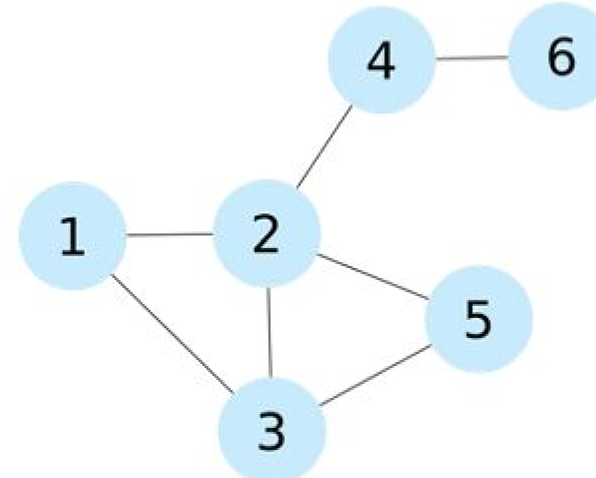
N6 → {N4}

- Pros:
  - Getting the neighbours of a node
  - Cheap addition of nodes
  - More compact representation of sparse graphs
- Cons:
  - Checking if there is an edge between two nodes
    - Optimization: sorted lists => logarithmic scan, but also logarithmic insertion

Data Structure: Incidence Matrix

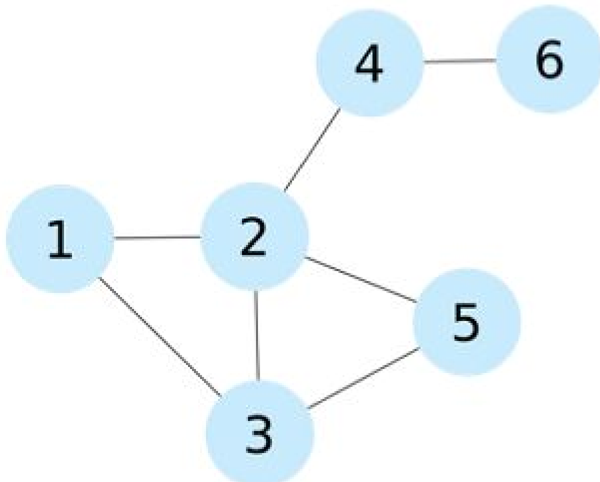


- Two-dimensional Boolean matrix of *n* rows and *m* columns
  - A column represents an edge
    - Nodes that are connected by a certain edge
  - A row represents a node
    - All edges that are connected to the node



$$\begin{pmatrix} 1 & 1 & 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 1 & 1 & 1 & 0 & 0 \\ 0 & 1 & 1 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 1 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{pmatrix}$$

Incidence Matrix: Properties

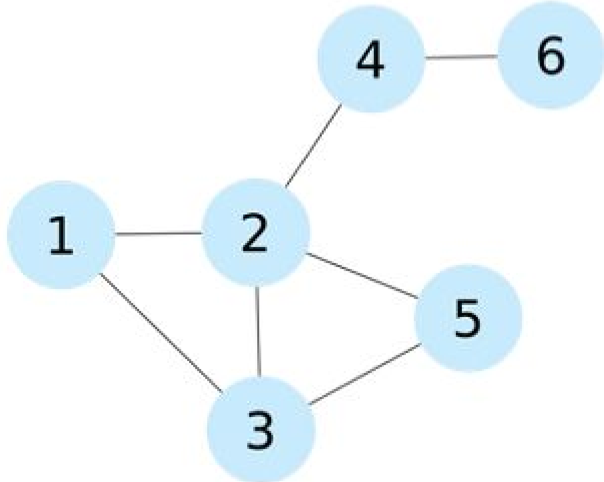


$$\begin{pmatrix} 1 & 1 & 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 1 & 1 & 1 & 0 & 0 \\ 0 & 1 & 1 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 1 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{pmatrix}$$

- Pros:
  - Representation of hypergraphs
    - where one edge connects an arbitrary number of nodes
- Cons:
  - Requires $n \times m$ bits (for most graphs $m \ggg n$)

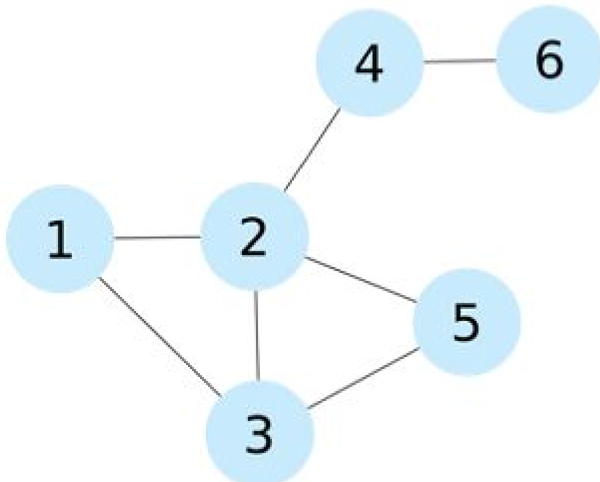
Data Structure: Laplacian Matrix

- Two-dimensional array of $n \times n$ integers
  - Similar structure as adjacency matrix
  - Diagonal of the Laplacian matrix indicates the degree of the node
  - The rest of positions are set to *-1* if the two vertices are connected, *0* otherwise



$$\begin{pmatrix} 2 & -1 & -1 & 0 & 0 & 0 \\ -1 & 4 & -1 & -1 & -1 & 0 \\ -1 & -1 & 3 & 0 & -1 & 0 \\ 0 & -1 & 0 & 2 & 0 & -1 \\ 0 & -1 & -1 & 0 & 2 & 0 \\ 0 & 0 & 0 & -1 & 0 & 1 \end{pmatrix}$$

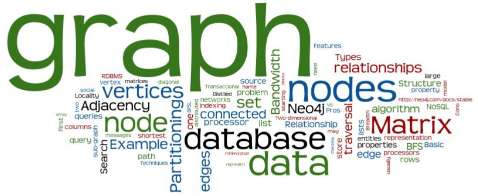
Laplacian Matrix: Properties



$$\begin{pmatrix} 2 & -1 & -1 & 0 & 0 & 0 \\ -1 & 4 & -1 & -1 & -1 & 0 \\ -1 & -1 & 3 & 0 & -1 & 0 \\ 0 & -1 & 0 & 2 & 0 & -1 \\ 0 & -1 & -1 & 0 & 2 & 0 \\ 0 & 0 & 0 & -1 & 0 & 1 \end{pmatrix}$$

All features of adjacency matrix

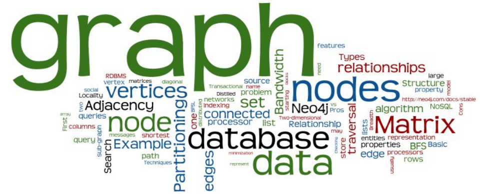
- Pros:
  - Analyzing the graph structure by means of spectral analysis
    - Calculating eigenvalues of the matrix



A Bit of a Theory

Selected graph algorithms

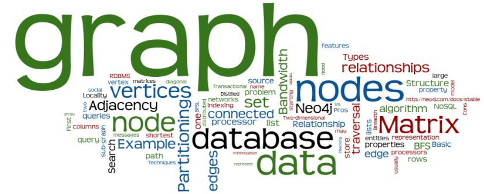
Basic Graph Algorithms



- Access all nodes:
  - Breadth-first Search (BFS)
  - Depth-first Search (DFS)
- Shortest path between two nodes
- Single-source shortest path problem
  - BFS (unweighted),
  - Dijkstra (nonnegative weights),
  - Bellman-Ford algorithm
- All-pairs shortest path problem
  - Floyd-Warshall algorithm

http://en.wikipedia.org/wiki/Shortest_path_problem

Improving Data Locality



- Performance of the read/write operations
  - Depends also on physical organization of the data
  - Objective: Achieve the best "data locality"

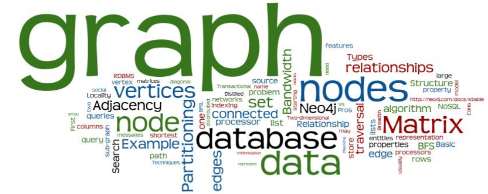
- Spatial locality:
  - if a data item has been accessed, the nearby data items are likely to be accessed in the following computations
    - e.g., during graph traversal

- Strategy:
  - in graph adjacency matrix representation, exchange rows and columns to improve the disk cache hit ratio
  - Specific methods: BFSL, Bandwidth of a Matrix, ...

Breadth First Search Layout (BFSL)

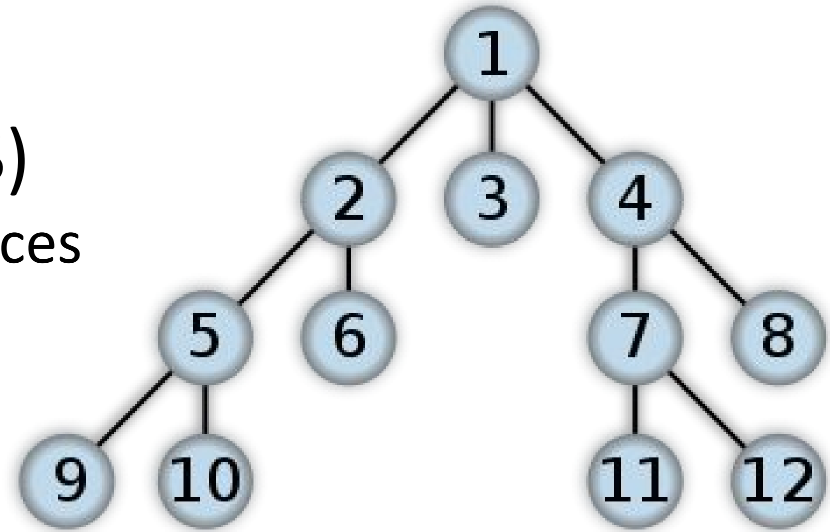
- Input: vertices of a graph
- Output: a permutation of the vertices
  - with better cache performance for graph traversals

- BFSL algorithm:
  1. Select a node (at random, the origin of the traversal)
  2. Traverse the graph using the BFS alg.
     - generating a list of vertex identifiers in the order they are visited
  3. Take the generated list as the new vertices permutation



Breadth First Search Layout (2)

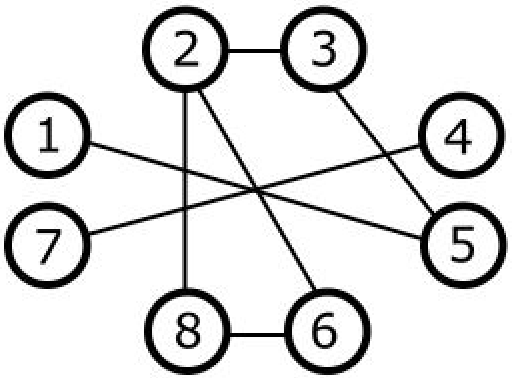
- Let us recall:
  Breadth First Search (BFS)
  - FIFO queue of frontier vertices

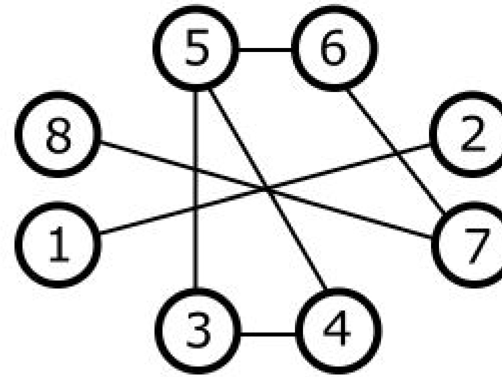


- Pros: optimal when starting from the same node
- Cons: starting from other nodes
  - The further, the worse

Matrix Bandwidth: Motivation

- Graph represented by adjacency matrix





$$
\begin{pmatrix}
1 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\
0 & 1 & 1 & 0 & 0 & 1 & 0 & 1 \\
0 & 1 & 1 & 0 & 1 & 0 & 0 & 0 \\
0 & 0 & 0 & 1 & 0 & 0 & 1 & 0 \\
1 & 0 & 1 & 0 & 1 & 0 & 0 & 0 \\
0 & 1 & 0 & 0 & 0 & 1 & 0 & 1 \\
0 & 0 & 0 & 1 & 0 & 0 & 1 & 0 \\
0 & 1 & 0 & 0 & 0 & 1 & 0 & 1
\end{pmatrix}
\rightarrow
\begin{pmatrix}
1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\
1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 1 & 1 & 1 & 0 & 0 & 0 \\
0 & 0 & 1 & 1 & 1 & 0 & 0 & 0 \\
0 & 0 & 1 & 1 & 1 & 1 & 0 & 0 \\
0 & 0 & 0 & 0 & 1 & 1 & 1 & 0 \\
0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 \\
0 & 0 & 0 & 0 & 0 & 0 & 1 & 1
\end{pmatrix}
$$

Matrix Bandwidth: Formalization

- The minimum bandwidth problem
  - Bandwidth of a row in a matrix = the maximum distance between nonzero elements, where one is left of the diagonal and the other is right of the diagonal
  - Bandwidth of a matrix = maximum bandwidth of its rows

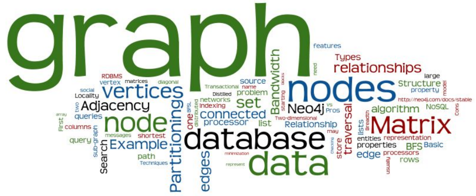
- Low bandwidth matrices are more cache friendly
  - Non zero elements (edges) clustered about the diagonal

- Bandwidth minimization problem: NP hard
  - For large matrices the solutions are only approximated

A Bit of a Theory

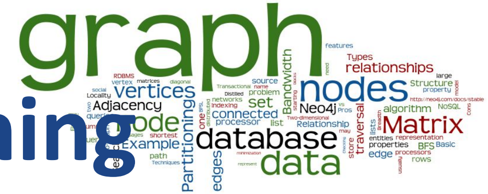
Graph partitioning

Graph Partitioning



- Some graphs are too large to be fully loaded into the main memory of a single computer
  - Usage of secondary storage degrades the performance
  - Scalable solution: distribute the graph on multiple nodes
- We need to partition the graph reasonably
  - Usually for a particular (set of) operation(s)
    - The shortest path, finding frequent patterns, BFS, spanning tree search
- This is difficult and graph DB are often centralized

Example: 1-Dimensional Partitioning



- Aim: partitioning the graph to solve BFS efficiently
  - Distributed into shared-nothing parallel system
  - Partitioning of the adjacency matrix

- 1D partitioning:
  - Matrix rows are randomly assigned to the *P* nodes (processors) in the system
  - Each vertex and the edges emanating from it are owned by one processor

	1	2	3	4	5	6	7	8	9	10	11	12
1	0	0	0	0	0	0	0	0	0	1	1	0
2	0	0	1	0	0	0	0	1	0	0	0	0
3	0	1	0	0	0	0	1	1	0	0	0	0
4	0	0	0	0	1	0	0	0	0	0	0	1
5	0	0	0	1	0	0	0	0	0	0	0	1
6	0	0	0	0	0	0	1	0	0	0	1	0
7	0	0	1	0	0	1	0	1	0	1	1	0
8	0	1	1	0	0	0	1	0	0	0	0	0
9	0	0	0	0	0	0	0	0	0	0	1	1
10	1	0	0	0	0	0	1	0	0	0	1	0
11	1	0	0	0	0	1	1	0	1	1	0	0
12	0	0	0	1	1	0	0	0	1	0	0	0

One-Dimensional Partitioning: BFS

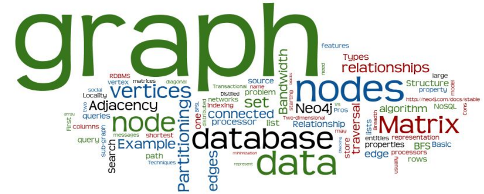
- BSF with 1D partitioning
  1. Each **processor** has a set of frontier vertices *F* (FIFO)
  2. The lists of neighbors of the vertices in *F* forms a set of neighbouring vertices *N*
     - Some owned by the current processor, some by others
  3. Messages are sent to all other processors… etc.

- 1D partitioning leads to high messaging
  - => 2D-partitioning of adjacency matrix
  - … lower messaging but still very demanding

Efficient sharding of a graph is very difficult

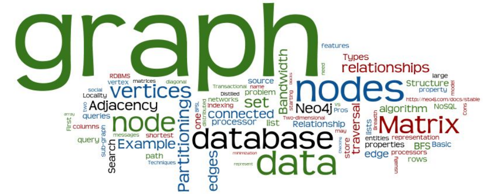
Graph Databases

Types of Graphs



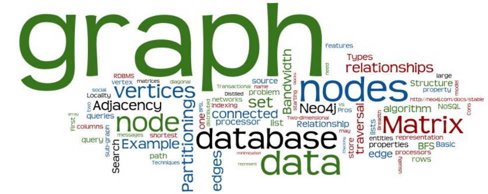
- Single-relational graphs
  - Edges are homogeneous in meaning
    - e.g., all edges represent friendship
- Multi-relational (property) graphs
  - Edges are typed or labeled
    - e.g., friendship, business, communication
  - Vertices and edges maintain a set of key/value pairs
    - Representation of non-graphical data (properties)
    - e.g., name of a vertex, the weight of an edge

Graph Databases



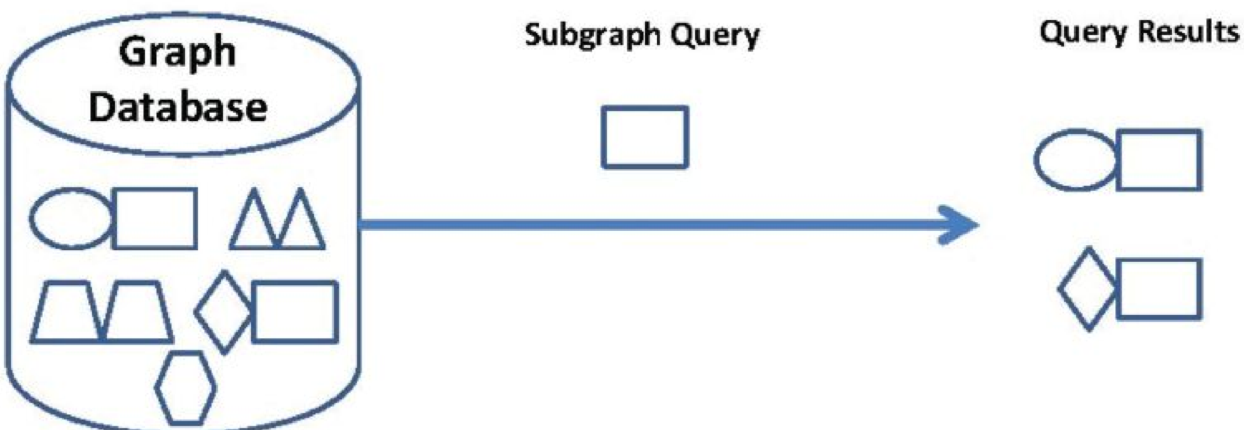
- A graph database = a set of graphs

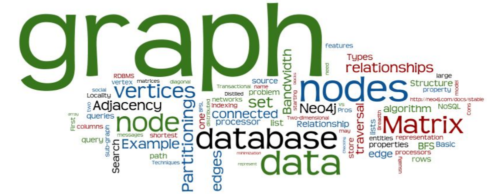
- Types of graph databases:
  - Transactional = large set of small graphs
    - e.g., chemical compounds, biological pathways, …
    - Searching for graphs that match the query
  - Non-transactional = few numbers of very large graphs
    - or one huge (not connected) graph
    - e.g., Web graph, social networks, …



Transactional DBs: Queries

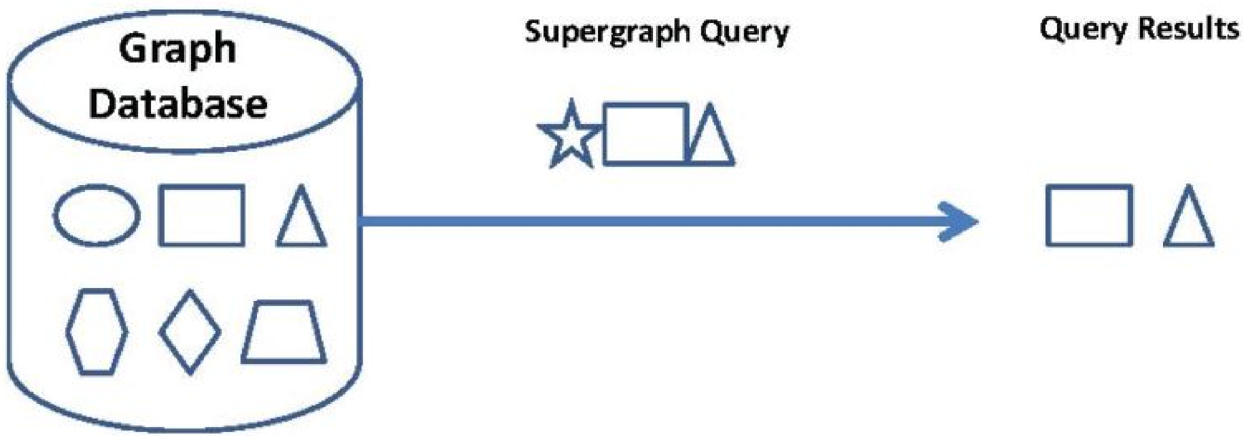
- Types of Queries
  - Subgraph queries
    - Searches for a specific pattern in the graph database
    - Query = a small graph or a graph, where some parts are uncertain
      - e.g., vertices with wildcard labels
    - More general type: allow sub-graph isomorphism





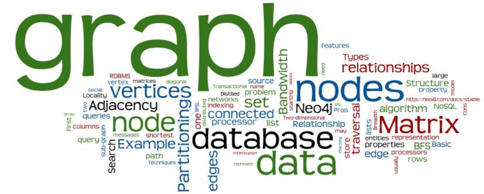
Transactional DBs: Queries (2)

- Super-graph queries
  - Search for the graph database members whose whole structure is contained in the input query



- Similarity (approximate matching) queries
  - Finds graphs which are similar to a given query graph
    - but not necessarily isomorphic
  - Key question: how to measure the similarity

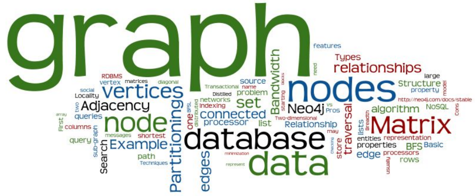
Indexing & Query Evaluation



- Extract certain characteristics from each graph
  - And index these characteristics for each $G_1, \ldots, G_n$

- Query evaluation in transactional graph DB
  1. Extraction of the characteristics from query graph q
  2. Filter the database (index) and identify a candidate set
     - Subset of the $G_1, \ldots, G_n$ graphs that should contain the answer
  3. Refinement - check all candidate graphs

Subgraph Query Processing



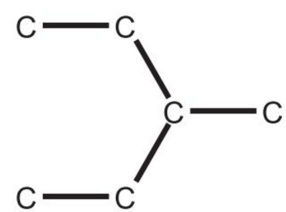
1. Mining-based Graph Indexing Techniques
   - Idea: if some features of query graph *q* do not exist in data graph *G*, then *G* cannot contain *q* as its subgraph
   - Apply graph-mining methods to extract some features (sub-structures) from the graph database members
     - e.g., frequent sub-trees, frequent sub-graphs
   - An inverted index is created for each feature

2. Non Mining-Based Graph Indexing Techniques
   - Indexing of the whole constructs of the graph database
     - Instead of indexing only some selected features

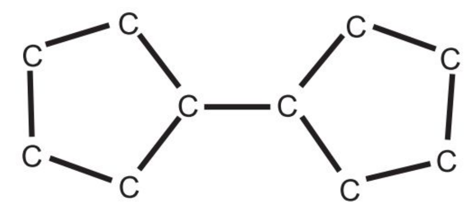
Mining-based Technique

- Example method: GIndex [2004]
  - Indexing “frequent discriminative graphs”
  - Build inverted index for selected discriminative subgraphs

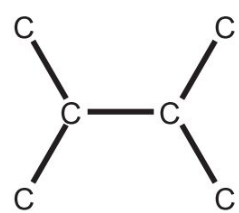
G_1



G_2



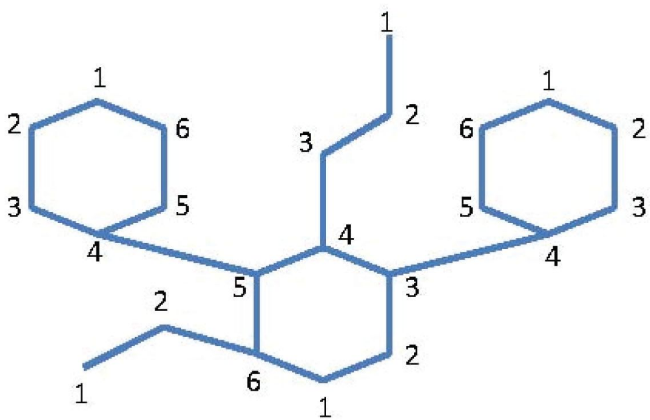
G_3

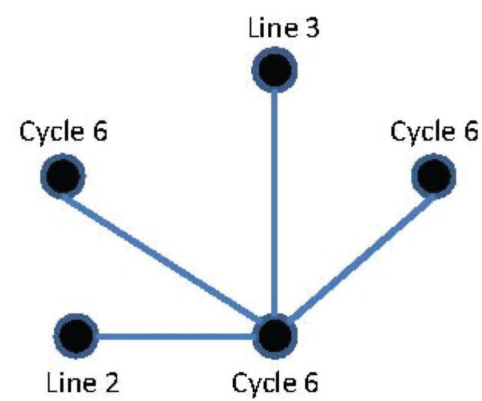


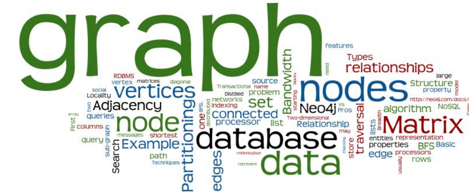
G_d

Non Mining-based Techniques

- **Example:** GString (2007)
  - Model the graphs in the context of organic chemistry using basic structures
    - Line = series of vertices connected end to end
    - Cycle = series of vertices that form a close loop
    - Star = core vertex directly connects to several vertices







Graph Databases

Non-transactional Databases

Non-transactional Databases

- A few very large graphs
    - e.g., Web graph, social networks, …
- Queries:
    - Nodes/edges with properties
    - Neighboring nodes/edges
    - Paths (all, shortest, etc.)
- Our example: Neo4j

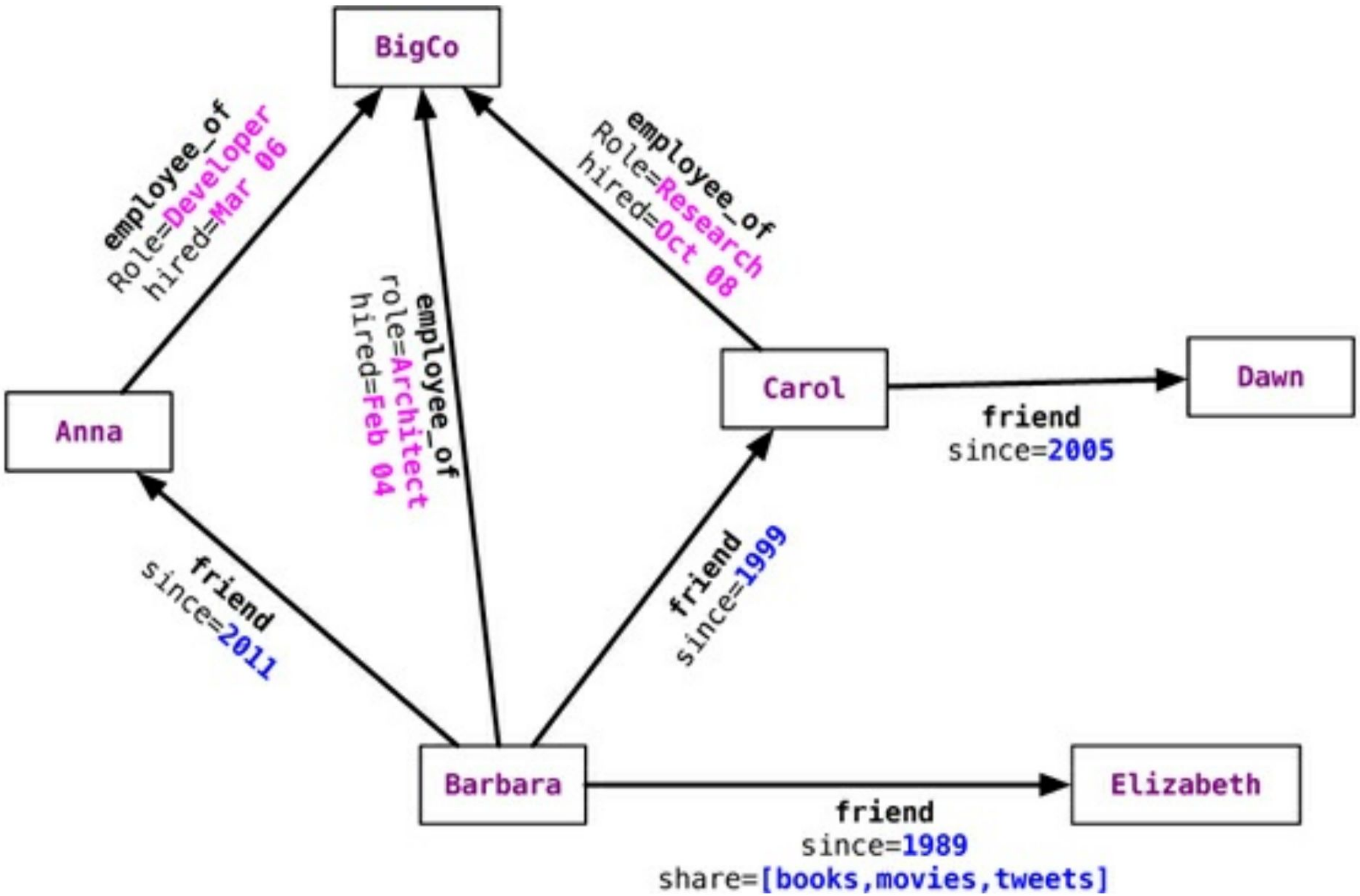
Basic Characteristics

- Different types of relationships between nodes
  - To represent relationships between domain entities
  - Or to model any kind of secondary relationships
    - Category, path, time-trees, spatial relationships, …

- No limit to the number and kind of relationships

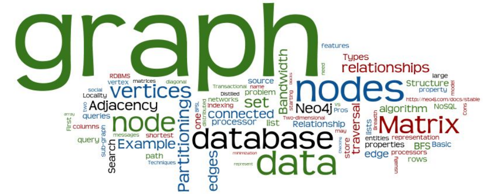
- Relationships have: type, start node, end node, own properties
  - e.g., "since when" did they become friends

Relationship Properties: Example

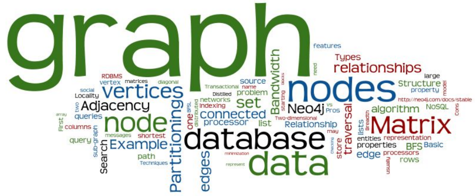


source: Sadalage & Fowler: NoSQL Distilled, 2012

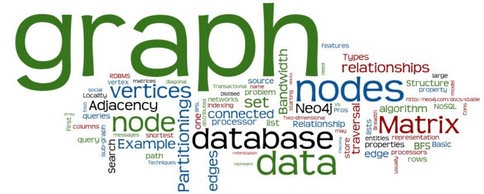
Graph DB vs. RDBMS



- RDBMS designed for a single type of relationship
  - “Who is my manager”
- Adding another relationship usually means a lot of schema changes
- In RDBMS we model the graph beforehand based on the traversal we want
  - If the traversal changes, the data will have to change
  - Graph DBs: the relationship is not calculated but persisted

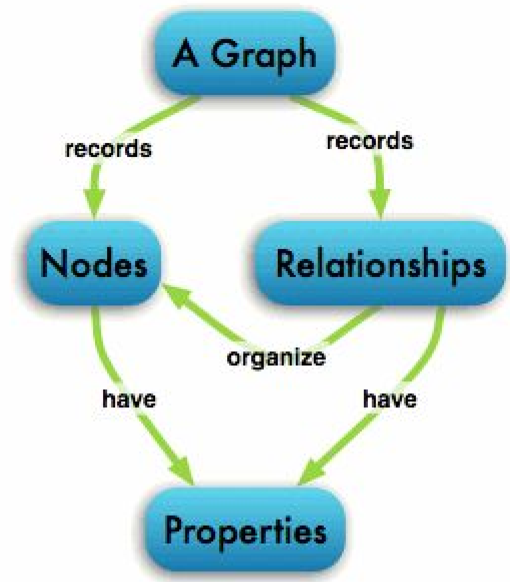


Neo4J: Basics & Concepts

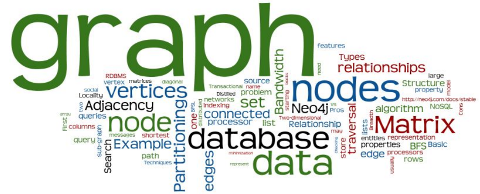


Neo4j: Basic Info

- Open source graph database
  - The most popular
- Initial release: 2007
- Written in: Java
- OS: cross-platform
- Stores data as nodes connected by directed, typed relationships
  - With properties on both
  - Called the “property graph”

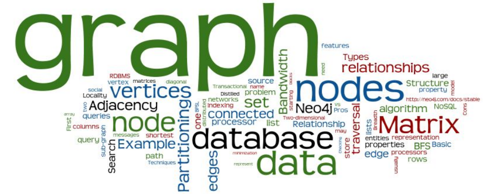


Neo4j: Basic Features

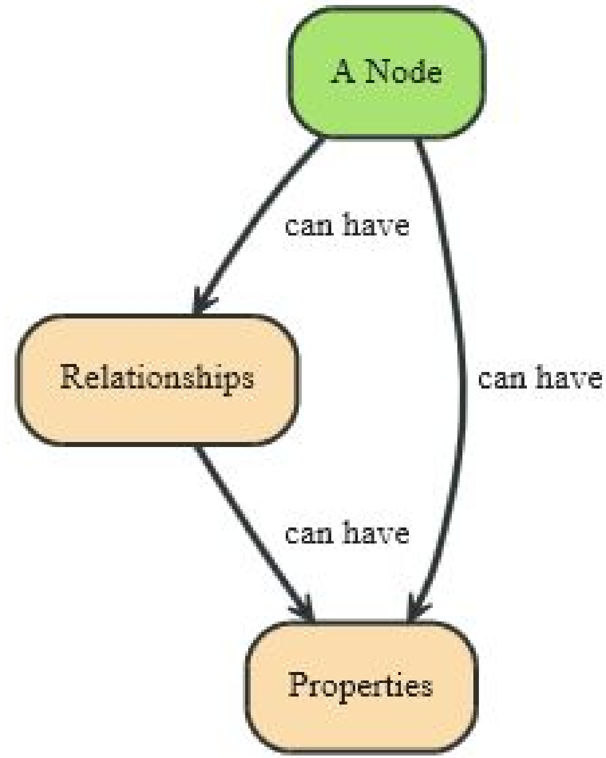


- reliable – with full ACID transactions
- durable and fast – disk-based, native storage engine
- scalable – up to several billion nodes/relationships/properties
- highly-available – when distributed (replicated)
- expressive – powerful, human readable graph query language
- fast – powerful traversal framework
- embeddable - in Java program
- simple – accessible by REST interface & Java API

Neo4j: Data Model



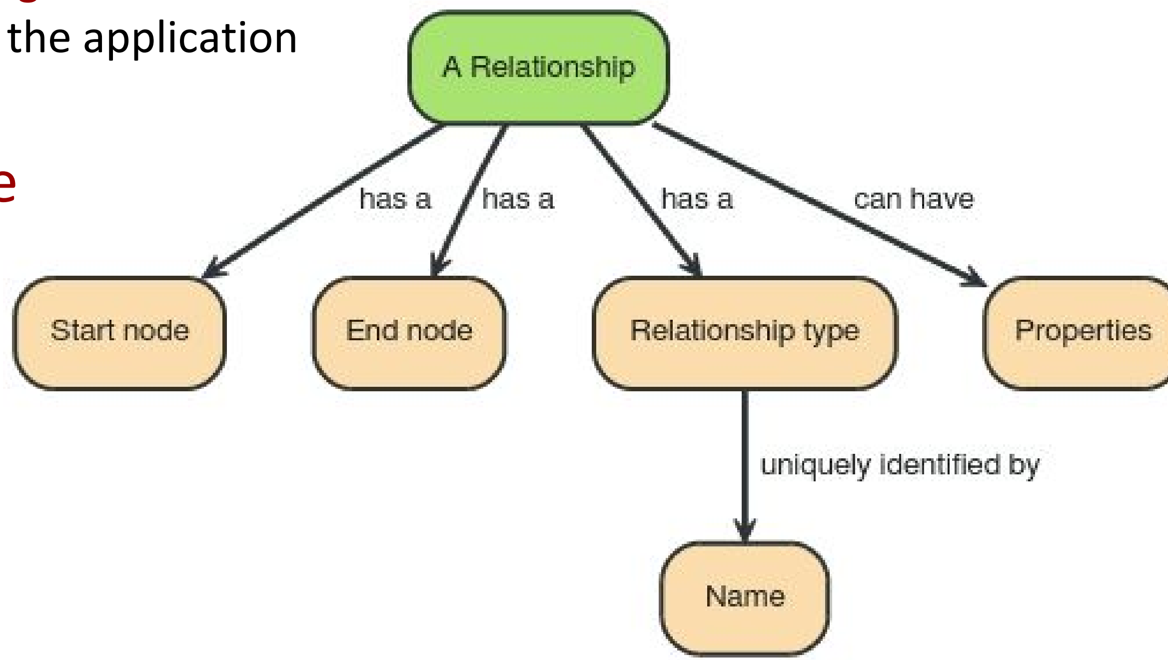
- Fundamental units: nodes + relationships
- Both can contain properties
  - Key-value pairs
  - Value can be of primitive type or an array of primitive type
  - null is not a valid property value
    - nulls can be modelled by the absence of a key



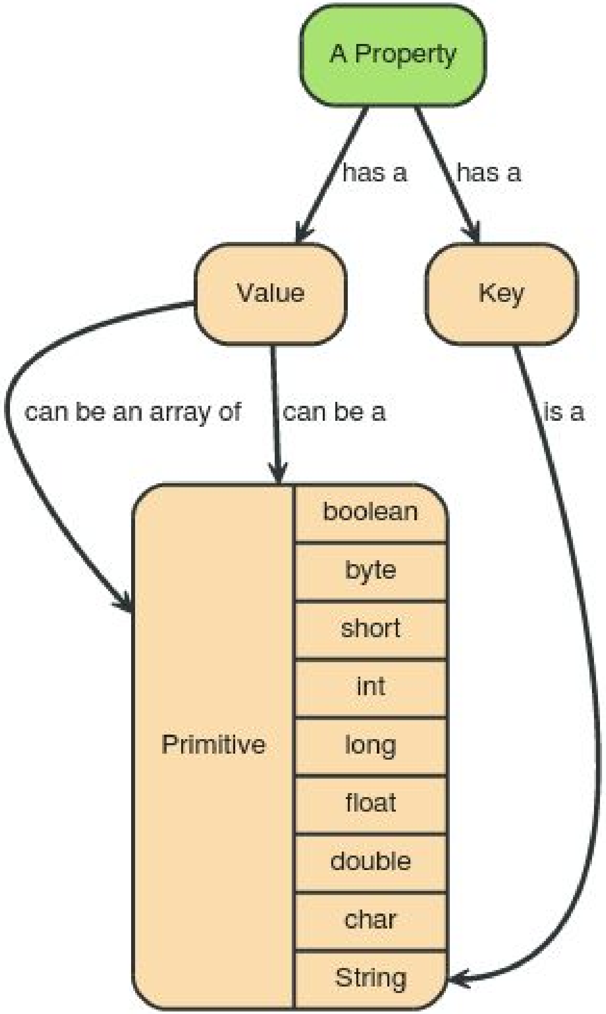
http://db-engines.com/en/system/Neo4j

Data Model: Relationships

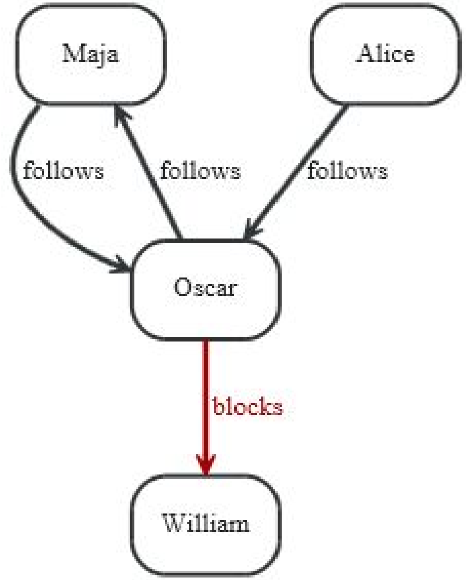
- Directed relationships (edges)
  - Incoming and outgoing edge
    - Equally efficient traversal in both directions
    - Direction can be ignored
      if not needed by the application
  - Always a start
    and an end node
    - Can be recursive



Data Model: Properties

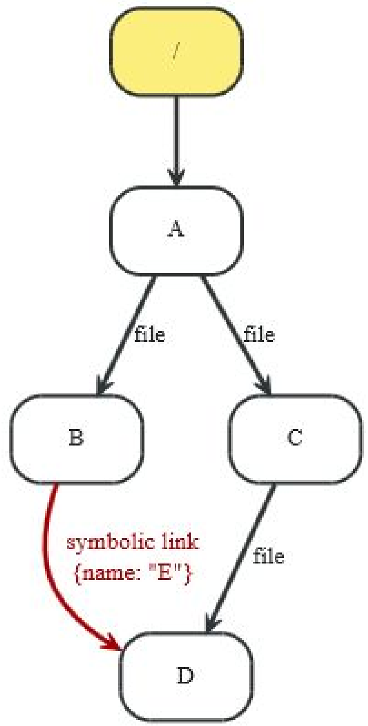


Type	Description
boolean	true/false
byte	8-bit integer
short	16-bit integer
int	32-bit integer
long	64-bit integer
float	32-bit IEEE 754 floating-point number
double	64-bit IEEE 754 floating-point number
char	16-bit unsigned integers representing Unicode characters
String	sequence of Unicode characters

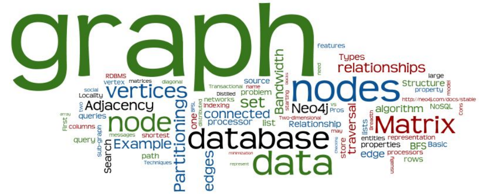


What	How
get who a person follows	outgoing *follows* relationships, depth one
get the followers of a person	incoming *follows* relationships, depth one
get who a person blocks	outgoing *blocks* relationships, depth one

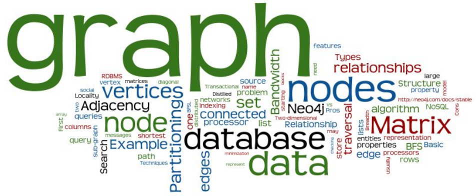
What	How
get the full path of a file	incoming *file* relationships
get all paths for a file	incoming *file* and *symbolic link* relationships
get all files in a directory	outgoing *file* and *symbolic link* relationships, depth one
get all files in a directory, excluding symbolic links	outgoing *file* relationships, depth one
get all files in a directory, recursively	outgoing *file* and *symbolic link* relationships



Access to Neo4j

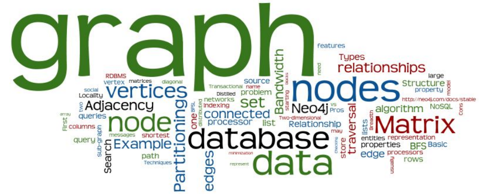


- Embedded database in Java system
- Language-specific connectors
  - Libraries to connect to a running Neo4j server
- Cypher query language
  - Standard language to query graph data
- HTTP REST API
- Gremlin graph traversal language (plugin)
- etc.



Neo4J: Native Java API & Graph Traversal

Native Java Interface: Example

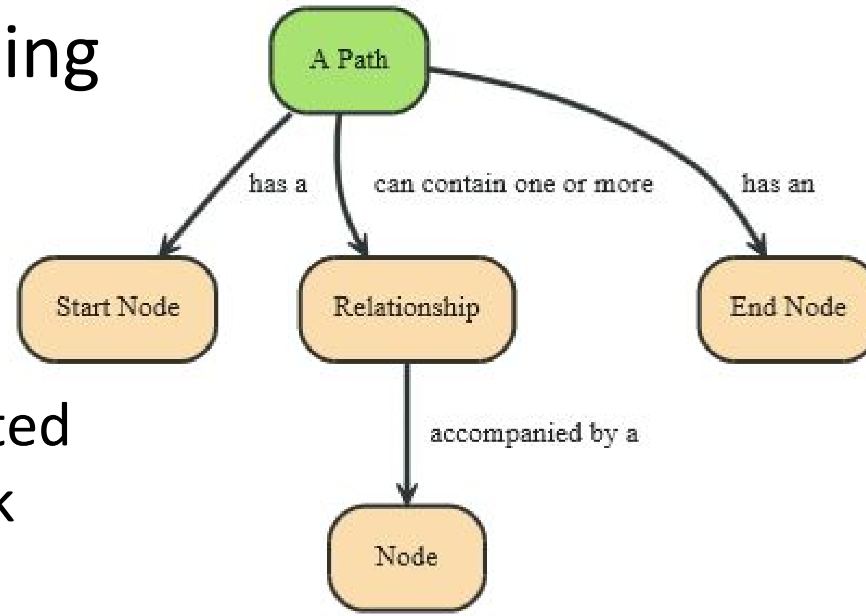



```
Node irena = graphDb.createNode();
irena.setProperty("name", "Irena");
Node jirka = graphDb.createNode();
jirka.setProperty("name", "Jirka");

Relationship i2j = irena.createRelationshipTo(jirka, FRIEND);
Relationship j2i = jirka.createRelationshipTo(irena, FRIEND);

i2j.setProperty("quality", "a good one");
j2i.setProperty("since", 2003);
```

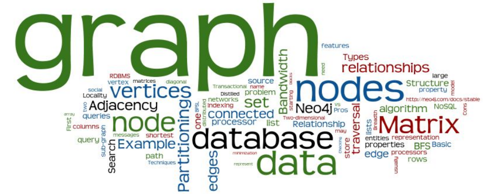
- Undirected edge:
  - Relationship between the nodes in both directions
  - **INCOMING** and **OUTGOING** relationships from a node

Data Model: Traversal + Path

- Path = one or more nodes + connecting relationships
  - Typically retrieved as a result of a query or a traversal
- Traversing a graph = visiting its nodes, following relationships according to some rules
  - Typically, a subgraph is visited
  - Neo4j: Traversal framework + Java API, Cypher, Gremlin



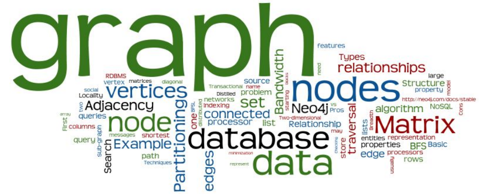
Traversal Framework



- A traversal is influenced by
  - Starting node(s) where the traversal will begin
  - Expanders – define what to traverse
    - i.e., relationship direction and type
  - Order – depth-first / breadth-first
  - Uniqueness – visit nodes (relationships, paths) only once
  - Evaluator – what to return and whether to stop or continue traversal beyond a current position

Traversal = TraversalDescription + starting node(s)

Traversal Framework – Java API



- `org.neo4j...TraversalDescription`
  - The main interface for defining traversals
    - Can specify branch ordering `breadthFirst()` / `depthFirst()`

- `.relationships()`
  - Adds the relationship type to traverse
    - e.g., traverse only edge types: FRIEND, RELATIVE
    - Empty (default) = traverse all relationships
  - Can also specify direction
    - `Direction.BOTH`
    - `Direction.INCOMING`
    - `Direction.OUTGOING`

Traversal Framework – Java API (2)

- `org.neo4j...Evaluator`
  - Used for deciding at each node: should the traversal continue, and should the node be included in the result
    - `INCLUDE_AND_CONTINUE`: Include this node in the result and continue the traversal
    - `INCLUDE_AND_PRUNE`: Include this node, do not continue traversal
    - `EXCLUDE_AND_CONTINUE`: Exclude this node, but continue traversal
    - `EXCLUDE_AND_PRUNE`: Exclude this node and do not continue
  - Pre-defined evaluators:
    - `Evaluators.toDepth(int depth)` / `Evaluators.fromDepth(int depth)`,
    - `Evaluators.excludeStartPosition()`
    - ...

Traversal Framework – Java API (3)

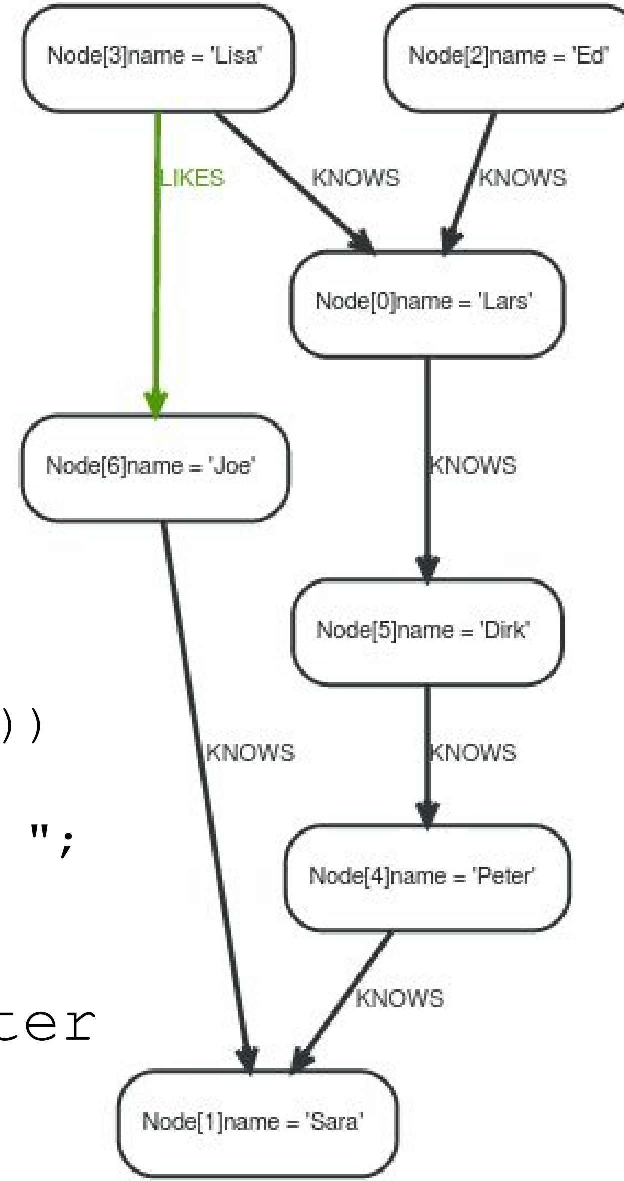
- `org.neo4j...Uniqueness`
  - Can be supplied to the TraversalDescription
  - Indicates under what circumstances a traversal may revisit the same position in the graph

- `Traverser`
  - Starts actual traversal given a TraversalDescription and starting node(s)
  - Returns an iterator over "steps" in the traversal
    - Steps can be: Path (default), Node, Relationship
  - The graph is actually traversed "lazily" (on request)

Example of Traversal

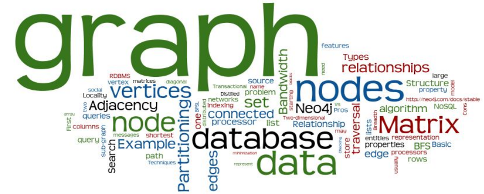
```
TraversalDescription desc =
  db.traversalDescription()
    .depthFirst()
    .relationships( Rels.KNOWS,
                    Direction.BOTH )
    .evaluator(Evaluators.toDepth(3));

// node is 'Ed' (Node[2])
for (Node n : desc.traverse(node).nodes())
{
    output += n.getProperty("name") + ", ";
}
```

Output: `Ed, Lars, Lisa, Dirk, Peter`

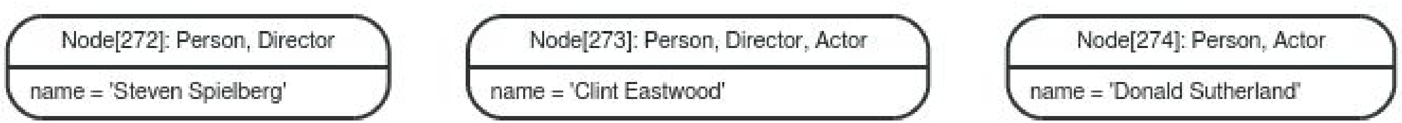


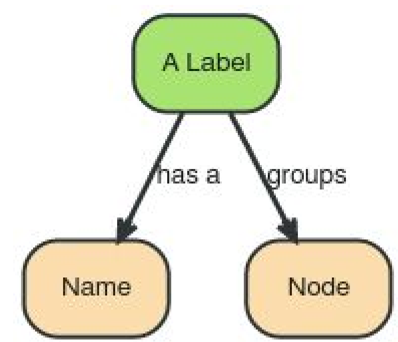
http://neo4j.com/docs/stable/tutorial-traversal-java-api.html

Access to Nodes



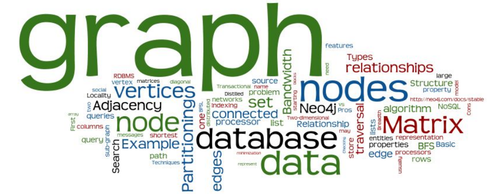
- How to get to the starting node(s) before traversal
  1. Using internal identifiers (unique generated IDs)
     - not recommended because Neo4j does reuse freed IDs
  2. Using specified properties
     - one of the properties is typically "ID" (natural user-specified ID)
     - recommended, properties can be indexed
       - automatic indexes
  3. Using "labels"
     - group nodes into "subsets" (named graph)
     - a node can have more than one label
       - belong to more subsets





Neo4J: Cypher Language

Cypher Language



- Neo4j graph query language
  - For querying and updating
- Declarative – we say what we want
  - Not how to get it
  - Not necessary to express traversals
- Human-readable
- Inspired by SQL and SPARQL
- Still growing = syntax changes are often

http://neo4j.com/docs/stable/cypher-query-lang.html

Cypher: Clauses

- **MATCH**: The graph pattern to match
- **WHERE**: Filtering criteria
- **RETURN**: What to return
- **CREATE**: Creates nodes and relationships.
- **DELETE**: Remove nodes, relationships, properties
- **SET**: Set values to properties
- **WITH**: Divides a query into multiple parts
- **START**: Starting points in the graph
  - by explicit index lookups or by node IDs (both deprecated)

Cypher: Creating Nodes (Examples)

CREATE n;

*(create a node, assign to var **n**)*

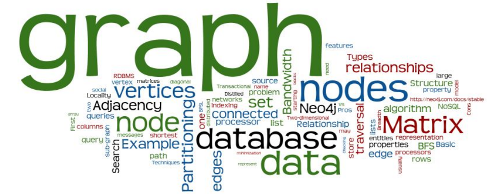
```
Created 1 node, returned 0 rows
```

CREATE (a: Person {name : 'David'})
RETURN a;

(create a node with label 'Person' and 'name' property 'David')

```
Created 1 node, set 1 property, returned 1 row
```

Cypher: Creating Relationships



START a=node(361), b=node(362)
CREATE a-[r:Friend]->b
RETURN r ;

(create relations Friend *between nodes with IDs 1 and 2)*

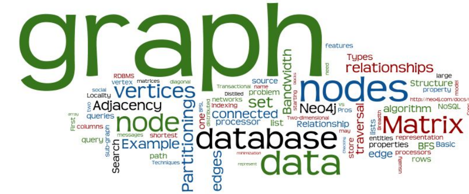
```
Created 1 relationship, returned 1 row
```

START a=node(1), b=node(2)
CREATE a-[r:Friend {name : a.name + '->' + b.name }]->b
RETURN r

(set property 'name' of the relationship)

```
Created 1 node, set 1 property, returned 1 row
```

Cypher: Queries




```
MATCH (p: Person)
WHERE p.age > 18 AND p.age < 30
RETURN p.name
```

(return names of all adult people under 30)

```
MATCH (user: Person {name: 'Andres'})-[:Friend]->(follower)
RETURN user.name, follower.name
```

(find all ‘Friends’ of 'Andres')

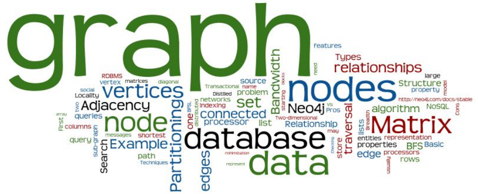
Cypher: Queries (2)

```
MATCH (andres: Person {name: 'Andres'})-[*1..3]-(node)
RETURN andres, node ;
```

(find all 'nodes' within three hops from 'Andres')

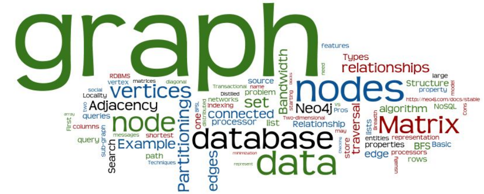
```
MATCH p=shortestPath(
  (andres:Person {name: 'Andres'})-[*]-(david {name:'David'})
)
RETURN p ;
```

(find the shortest connection between 'Andres' and 'David')



Neo4J: Behind the Scene

Neo4j Internals: Indexes




```
CREATE INDEX ON :Person(name);
```

(Create index on property name from label Person)

```
Indexes added: 1
```

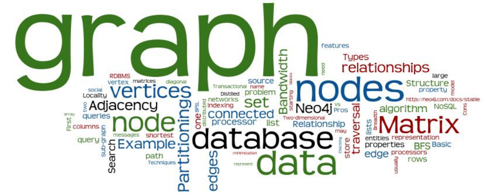
- Since Neo4j v. 2, indexes are used automatically
- Can be specified explicitly (which index to use)

```
MATCH (n:Person)
USING INDEX n:Person(surname)
WHERE n.surname = 'Taylor'
RETURN n
```

Neo4j Internals: Transactions

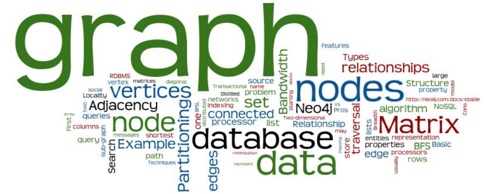
- Transactions in Neo4j
  - Support for ACID properties
  - All write operations must be performed in a transaction
  - Transaction isolation level: Read committed
    - Operation can see the last committed value
    - Reads do not block or take any locks
    - If the same row is retrieved twice within a transaction, the values in the row CAN differ
  - Higher level of isolation can be achieved
    - By explicit acquiring the read locks

Neo4j Internals: High Availability



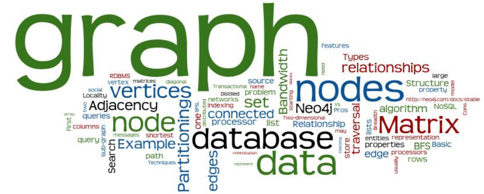
- Master-slave replication
  - Several Neo4j slave databases can be configured to be exact replicas of a single Neo4j master database
- Speed-up of read operations
  - A horizontally scaling read-mostly architecture
  - Enables to handle more read load than a single node
- Fault-tolerance
  - In case a node becomes unavailable
- Transactions are still atomic, consistent and durable, but eventually propagated to the slaves

Graph DBs: Suitable Use Cases



- Connected Data
  - Social networks
  - Any link-rich domain is well suited for graph databases
- Routing, Dispatch, and Location-Based Services
  - Node = location or address that has a delivery
  - Graph = nodes where a delivery has to be made
  - Relationships = distance
- Recommendation Engines
  - “your friends also bought this product”
  - “when buying this item, these others are usually bought”

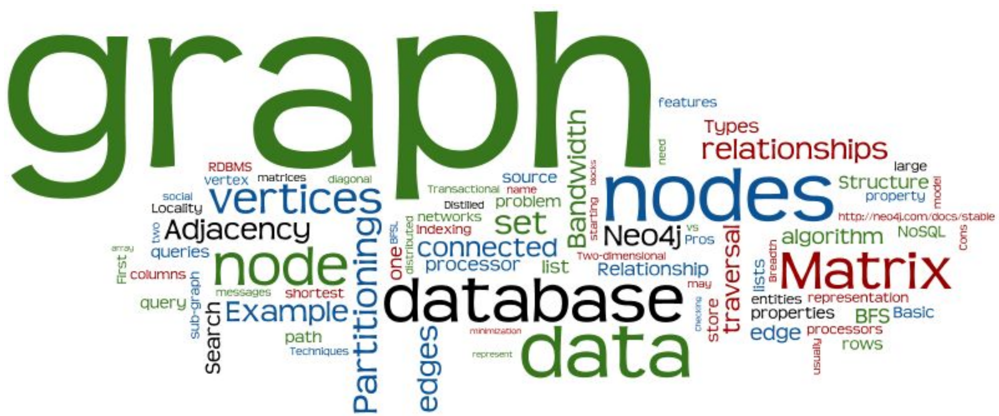
Graph DBs: When Not to Use



- If we want to update all or a subset of entities
  - Changing a property on many nodes is not straightforward
    - e.g., analytics solution where all entities may need to be updated with a changed property
- Some graph databases may be unable to handle lots of data
  - Distribution of a graph is difficult

Questions?

Please, any questions? Good question is a gift...



References

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