Constraint Programming: from LP to CLP

Eric Monfroy

IRIN, Université de Nantes

From Logic Programming to Constraint Logic Programming

Objectives

- historical integration of CP in LP
- reminder of LP
- main difference between LP and CLP (execution)

Brief historical review

- 1963. Sketchpad: introduction of CP techniques (Interactive drawing.)
- 1975. First formalism for CP techniques: scene labeling (recognition of 3D objects from 2D drawings)
- 1985. Discovering that logic programming is a special instance of constraint programming :

unification = resolution of constraints over trees

Moreover: Prolog is a relationnal and declarative language, and it offers features to explore the search space (backtracking)

- ⇒ Constraint Logic Programming (CLP)
- Constraints supply a relationnal aspect to Prolog arithmetic (vs. the is/2 predicate)

Prolog syntax reminder (1)

- Variables: starts with an upper-case letter or underscore
 X, _e34, Variable3
- Constants: starts with a lower-case letter a, 'character quotes', pi
- Structure: date(sunday, X, year(1999))
- Term : variable, constant, or structure (= data structure of the program)
- Atom : expression of the form p(t1,...,tn) where p is a predicate symbol and the ti's are terms
- Fact: expression of the form p(t1,...,tn).
 father(john,X)

Prolog syntax reminder (2)

- Paule : expression of the form :
 p(t,...,t) :- p(t,...,t), ..., p(t,...,t).

 Example :
 menu(E,P,D) :- starter(E), main_dish(P), dessert(D).
- Head : left-hand side of the rule
- Body : right-hand side of the rule
- Clause : rule or fact
- Query : clause without head

Prolog syntax reminder (3)

Predicate: set of clauses whose head have the same name and the same arity

Example:

```
dog(rex).
dog(X) :- pet(X), bark(X).
dog(X) :- bite(X).
```

- "," : conjunction
- ";": disjunction
- a fact describe a basic truth : « rex is a dog »
- meaning of the 2nd clause : « X is a dog if X is a pet and X barks ».
- each clause of a predicate is an alternative (implicit "or")
- the name of a variable is local to the clause

Prolog execution

- procedural view of clauses, where call with parameter passing is replaced by call with variable unification.
- unification of terms: minimal substitution of variables that make the 2 terms equal.
- no description of the operationnal aspect: the order to consider clauses and goals is arbitrary (in theory).
- use of backtracking to explore all the alternatives

LP versus CLP (1)

Prolog approach :

```
p(X,Y,Z) := Z \text{ is } X+Y.
:- p(3,4,Z).
Z = 7
:- p(X, 4, 7).
```

INSTANTIATION ERROR

CLP approach :

$$p(X,Y,Z) :- Z #= X+Y.$$
 $:- p(3,4,Z).$
 $Z = 7$
 $:- p(X,4,7).$
 $X = 3.$

LP versus CLP (2)

Prolog arithmetic is not relationnal

 \Rightarrow the only possible approach is *generate and test*:

```
solution(X,Y,Z):- p(X), p(Y), p(Z), test(X,Y,Z).

p(11). p(3). p(7). p(16). p(15). p(14).

test(X,Y,Z):- Y is X+1, Z is Y+1.
```

458 steps to get the first solution

:- solution (X,Y,Z).

LP versus CLP (3)

CLP arithmetic is relationnal

⇒ possible approach : *constrain and generate*

```
solution (X, Y, Z) := test(X, Y, Z), p(X), p(Y), p(Z).

p(11). p(3). p(7). p(16). p(15). p(14).

test(X, Y, Z) := Y #= X+1, Z #= Y+1.
```

```
:- solution (X, Y, Z).
```

11 steps to get the first solution